

# The Mad Hatters

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# Recreational Puzzles

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- ▶ Winkler, 2004 Mathematical Puzzles: A Connoisseur's Collection

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- ▶ What is the optimal strategy?

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- ▶ Players who guess wrong are eaten, those who guess right get shown the way off the island.



# Example-Example-People in a Line

□	■	□	□	■
Back				Front

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- ▶ This does at least identify what's hard here: We can't share information!

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- ▶ Yes! Both players assume that they have the same colour hat (50-50 chance).
- ▶ Win half the time!!!

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- ▶ First less likely hat colour is best.

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- ▶ Four things Player 1 could see, so 16 possible strategies for him.
- ▶ 256 group strategies.

# Two hats.

Two hats	■ ■	■ □	□ ■	□ □
■ ■				
■ □				
□ ■				
□ □				

# Two hats.

Two hats	■ ■	■ □	□ ■	□ □
■ ■	lose	lose	lose	lose
■ □	lose			
□ ■	lose			
□ □	lose			win

# Two hats.

Two hats	■ ■	■ □	□ ■	□ □
■ ■	lose	lose	lose	lose
■ □	lose			share
□ ■	lose			share
□ □	lose	share	share	win



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# Three hats.

Three hats	$\emptyset$	{1}	{2}	{1, 2}	{3}	{1, 3}	{2, 3}	{1, 2, 3}
Picture	■■■	□■■	■□■	□□■	■■■□	□■□	■□□	□□□
Choice	any	1	3	1	2	2	3	any

Table : Optimal strategy on 3 hats

For convenience

White hats	$\emptyset$	{1}	{2}	{1, 2}	{3}	{1, 3}	{2, 3}	{1, 2, 3}
Picture	■■■	□■■	■□■	□□■	■■■□	□■□	■□□	□□□
Choice	1	1	3	1	2	2	3	1

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# Three hats-example.

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□	■
■	□
□	□
Player 1	Player 2

Table : Player 1 chooses hat 1, Player 2 hat 2. They win

# Three hats.

White hats	■■■	■■□	■□■	■□□	□■■	□■□	□□■	□□□
■■■								
■■□			w	w				
■□■		w				w		
■□□		w		w		w		
□■■					w		w	w
□■□		w	w				w	w
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- ▶ Wait a lot of generations

- ▶ Rerun a lot

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- ▶ Took there "natural" infinite analogs.

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- ▶  $S_3$  Dual of  $S_1$ . Toggle all colours and play  $S_1$

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- ▶  $S_2$  was easier because fewer interactions.
- ▶  $S_3$  computed as the dual of  $S_1$ .

$$\begin{aligned}V_{S^d}(p) &= \mathbb{P}(A_{S^d}^{W,W}(p)) \\ &= \mathbb{P}(A_S^{B,B}(q)) \\ &= p - \mathbb{P}(A_S^{B,W}(q)) \\ &= p - (q - \mathbb{P}(A_S^{W,W}(q))) \\ &= p - q + \mathbb{P}(A_S^{W,W}(q)) \\ &= 2p - 1 + V_S(q)\end{aligned}$$



For our game with probability  $p$  of each hat being white, this strategy gives the following lower bound on  $V(p)$ :

1.  $\frac{p(1 + p + p^2 + 3p^3 - 3p^4 + p^5)}{(1 + p)(2 - p)(1 + p^2)} \leq V(p)$  for  $p \leq \frac{1}{2}$ ;
2.  $\frac{p(1 + 5p - 10p^2 + 10p^3 - 5p^4 + p^5)}{(2 - 2p + p^2)(1 + p)(2 - p)}$  for  $\frac{1}{2} \leq p$ .

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- ▶ First bound better for  $p < 1/2$ , second bound better for  $p > 1/2$ .
- ▶ Lowest terms of  $a$  and  $b$  is strongest, works best for  $\binom{b}{a}$  small.

- ▶ Multiple hat colours

# Future Work

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- ▶ Both multiple colours and multiple players.

# Thanks

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